

A simple construction for perfect factors in the de Bruijn graph

Chris Mitchell

c.mitchell@rhul.ac.uk

<http://www.isg.rhul.ac.uk/~cjm>

Agenda

1. de Bruijn sequences and arrays
2. Perfect factors
3. A new construction method
4. Further thoughts

De Bruijn sequences

- A (span v , c -ary) de Bruijn sequence is an infinite periodic sequence of symbols $\{0, 1, \dots, c-1\}$ (for some $c > 1$), with the property that every possible v -tuple of symbols occurs exactly once in a period.
- The period must clearly be equal to the number of c -ary v -tuples, i.e. c^v .

Examples

- [00011101] is a span 3, binary (2-ary) de Bruijn sequence (of length $2^3=8$).
- Here, as throughout, we write just one period of the bi-infinite sequence, and call it a **cycle** – we use square brackets for cycles.
- [001122021] is a span 2, 3-ary de Bruijn sequence (of length $3^2=9$).

Simple constructions I

- One very well-known method of generating binary de Bruijn sequences is a greedy algorithm, known as the ‘prefer one’ method.
- Start with the all zero v -tuple, and add one bit at a time, always adding a one **if possible**.
- E.g., for $v=4$:

[0000111101100101].

Simple constructions II

- Choose a prime power q .
- Take a maximum length v -stage shift register sequence over $GF(q)$ (an m -sequence).
- This is a q -ary periodic sequence of period $q^v - 1$, in which every q -ary v -tuple occurs except the all zero tuple.
- Take any one of the $(q-1)$ all zero $(v-1)$ -tuples in a period and insert a zero – this gives a span v q -ary de Bruijn sequence.

Pseudorandom sequences

- Equally, one can remove a zero from the unique all-zero v -tuple in any de Bruijn sequence to obtain a sequence which contains all possible c -ary v -tuples except the all-zero tuple.
- Such a sequence is known as a pseudorandom sequence.

Example

- [001122021] is a span 2, 3-ary de Bruijn sequence (of length $3^2=9$).
- [01122021] is a span 2, 3-ary pseudorandom sequence.
- [011220021] is another span 2, 3-ary de Bruijn sequence (of length $3^2=9$).

Existence

- De Bruijn (1946) and Good (1946) independently proved that de Bruijn sequences exist for every possible span v and every possible alphabet size $c > 1$.
- They also gave an explicit formula for the number of such sequences for every c and v .
- It was subsequently discovered that their existence had first been established by Flye-Sainte Marie in 1894.
- Since the 1940s a significant number of different construction methods have been devised.
- Fredricksen (1982) gave an extremely helpful summary of construction techniques (since 1982 more techniques found).

de Bruijn-Good graph I

- This directed graph (which we write as $G(c, v)$) has vertices the c -ary v -tuples.
- Put an edge from vertex $(a_0, a_1, \dots, a_{v-1})$ to vertex $(b_0, b_1, \dots, b_{v-1})$ if and only if $a_{i+1} = b_i$ (for $0 \leq i \leq v-2$).
- A **Hamiltonian cycle** in $G(c, v)$ then corresponds to a c -ary span v de Bruijn sequence.

de Bruijn-Good graph II

- We can also label edges of $G(c, v)$ with c -ary $(v+1)$ -tuples, i.e. so that the edge connecting $(a_0, a_1, \dots, a_{v-1})$ to $(a_1, a_2, \dots, a_{v-1}, b)$ is labelled $(a_0, a_1, \dots, a_{v-1}, b)$.
- It is then not hard to see that an **Eulerian cycle** in $G(c, v)$ corresponds to a c -ary span $v+1$ de Bruijn sequence.
- Since the in-degree of every vertex is equal to the out-degree ($=c$) such an Eulerian cycle always exists – hence establishing the existence of de Bruijn sequences.

de Bruijn arrays (perfect maps)

- de Bruijn arrays are 2-dimensional analogues of de Bruijn sequences.
- An $(m, n; u, v)_c$ -PM is a c -ary periodic array of period $m \times n$, in which every c -ary $u \times v$ sub-array occurs precisely once in a period (2-dimensional cycle).
- Hence, we must have $c^{uv} = mn$.
- Notion introduced by Reed and Stewart in 1962, who gave a $(4, 4; 2, 2)_2$ -PM.

Examples

- The Reed and Stewart example is as follows:

0	0	0	1
0	0	1	0
1	0	1	1
0	1	1	1

- It is straightforward to verify that every 2×2 binary array occurs once when this is regarded as the recurring periodic pattern in an infinite array (i.e. a 2-dimensional cycle).

Necessary conditions

- The obvious necessary condition for the existence of a $(m,n;u,v)_c$ -PM is $c^{uv}=mn$.
- For reasons it is simple to verify, we must also have:
 - i. $u=m=1$ or $1 \leq u < m$;
 - ii. $v=n=1$ or $1 \leq v < n$.
- These necessary conditions have been conjectured to be sufficient.

Existence

- The necessary conditions have been shown to be sufficient for the following cases:
 - $c=2$ (Paterson, 1994);
 - c a prime power (Paterson, 1996);
 - $u=v=2$ (Hurlbert, Mitchell and Paterson, 1996).

Applications

- A wide variety of applications have been proposed for de Bruijn sequences, including in:
 - cryptography;
 - position location.
- Position location/range finding applications have also been discussed for the two-dimensional arrays.

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Definition

- A perfect factor of a graph is a set of disjoint cycles of fixed length (n , say) which cover every edge.
- A perfect factor of the de Bruijn Graph $G(c, v-1)$ (which we call an $(n, v)_c$ -PF) can be thought of as a set of c^v/n periodic c -ary sequences (of period n), for which every c -ary v -tuple occurs precisely once in a period of just one of the sequences.

Examples

- If $n=c^v$, then an $(n, v)_c$ -PF is simply a span v , c -ary de Bruijn sequence.
- The set of two cycles: $\{[0,0,0,1], [1,1,1,0]\}$ forms a $(4,3)_2$ -PF.
- The set of three cycles: $\{[0,0,1], [1,1,2], [2,2,0]\}$ forms a $(3,2)_3$ -PF.
- The set of four cycles: $\{[0,0,3,3], [2,0,1,3], [1,1,2,2], [0,2,3,1]\}$ forms a $(4,2)_4$ -PF.

Necessary conditions

- We have the following trivial necessary conditions for the existence of an $(n, v)_c$ -PF:
 1. $n|c^v$;
 2. $v=n=1$ or $1 \leq v < n$.
- These necessary conditions have been conjectured to be sufficient.

Applications

- Perfect Factors are simply a special case of perfect maps, since any ordering of the c^v/n cycles as the columns of a c -ary periodic array will form a $(c^v/n, n; 1, v)_c$ -PM.
- [The converse is also trivially the case].
- Perfect Factors can also be used to help construct a much larger class of perfect maps (as we next see).

Etzion construction I

- Etzion (1988) showed how perfect factors can be used to construct perfect maps, generalising a construction of Ma (1984).
- We describe a special case of this construction proposed by Mitchell and Paterson in 1994.
- For simplicity we describe it for the binary case – however it works for arbitrary size alphabets.

Etzion construction II

- Let C_0, C_1, \dots, C_{n-1} be the cycles of an $(2^k, u)_2$ -PF [and hence $k \leq u$].
- Let (r_j) be $2^{k(v-1)}$ repetitions of a (2^{u-k}) -ary span v de Bruijn sequence [where $uv \geq 2k+1$].
- Let (s_j) be $2^{v(u-k)}$ repetitions of a 2^k -ary span $(v-1)$ pseudorandom sequence for which the first $v-2$ elements are all zeros, preceded by $2^{v(u-k)}$ zeros.
- Let (w_j) be defined so that $w_0=0$, $w_1=s_0$,
 $w_2=s_0+s_1$, $w_3=s_0+s_1+s_2$, ...

Etzion construction III

- Then define a $2^k \times 2^{uv-k}$ array made up of columns from the perfect factor.
- Specifically, let the i th column consist of cycle C_{r_i} cyclically shifted by w_i places.
- This is a $(2^k, 2^{uv-k}; u, v)_2$ -PM.
- Along with related constructions, this means that, if the perfect factor existence conjecture is positively resolved, the PM existence question will also be *mostly* resolved.

Existence I

- Etzion (1988) showed that $(2^k, v)_2$ -PFs exist if $k \leq v < 2^k$ [i.e. the necessary conditions are sufficient in $c=2$ case].
- Paterson (1994) showed the necessary conditions for an $(n, v)_c$ -PF are sufficient if c is a prime power.
- Mitchell (1994) showed the necessary conditions are sufficient for all allowable triples $(n, v)_c$ as long as there exists a prime p such that $p^\alpha | n$ and $p^\alpha > v$.
- Mitchell and Paterson (1998) observed that, to completely resolve the existence question, it is only necessary to establish the existence of an $(n, v)_c$ -PF for a 'square-free' c . This, in turn, means we only need to look at a finite number cases for each v .

Existence II

- The existence question has also been resolved for small values of v (the span).
- An $(n, v)_c$ -PF always exists if:
 - $v=2$ (Mitchell, 1994);
 - $v \leq 4$ (Mitchell, 1995);
 - $v \leq 6$ (Mitchell and Paterson, 1998).
- It was also established that if a $(10, 7)_{10}$ -PF and a $(10, 8)_{10}$ -PF could be constructed, then the existence conjecture would be resolved for $v \leq 8$.

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Objective

- The main goal is to give a new construction method.
- We first describe a very simple construction which forms the basis of the new method.
- This method first appears in the 1998 Mitchell-Paterson paper, but I have a feeling it was known before.

A simple construction method

- Suppose n, c are integers greater than 1 such that $n|c^{n-1}$.
- Consider the set of c -ary cycles of length n whose elements sum to a value congruent to 1 modulo c – since $n|c^{n-1}$, these cycles have period exactly c .
- If we regard cyclic shifted sequences as equivalent, we obtain an $(n, n-1)_c$ -PF.

Lempel homomorphism I

- We also need a homomorphism of the de Bruijn graph first given by Lempel (1970).
- The Lempel homomorphism D maps $G(c, v)$ to $G(c, v-1)$, and is defined by:
$$D(a_0, a_1, \dots, a_{v-1}) = (a_1 - a_0, a_2 - a_1, \dots, a_{v-1} - a_{v-2})$$
- D is a graph homomorphism – it is simple check that if there is an edge from \mathbf{a} to \mathbf{b} , then there is an edge from $D(\mathbf{a})$ to $D(\mathbf{b})$.

Lempel homomorphism II

- This homomorphism turns out to be an incredibly useful tool in the study of de Bruijn sequences and related structures.
- It is analogous to differentiation, and has many related properties.

Lempel homomorphism III

- We can apply D to periodic sequences, as well as just to tuples.
- If the sequence (s_i) has period u , then $(D(s_i))$ will have period dividing u .
- Also, the mod c sum of u consecutive elements of $(D(s_i))$ is always zero.

The inverse homomorphism

- If $\mathbf{s}=[s_0, s_1, \dots, s_{n-1}]$ is a cycle of weight w (reduced mod c) then we define the pre-image of \mathbf{s} , written $D^{-1}(\mathbf{s})$, to be the set of cycles:

$$\{ [t, t+s_0, t+s_0+s_1, \dots, t+(s_0+s_1+\dots+s_{n-2}), t+w, t+w+s_0, \dots] \}.$$

- This set has size (w,c) , and the cycles have period $nc/(w,c)$.
- Hence, if $w \bmod c = 0$, then $|D^{-1}(\mathbf{s})|=c$, and the cycles have period n .

The construction

- Suppose $c|n$ and c is odd.
- Suppose S is an $(n, n-1)_c$ -PF constructed using the simple construction method outlined previously.
- Then $D(S)$ is an $(n, n-2)_c$ -PF.

Example

- The set S of three cycles: $\{[0,0,1], [1,1,2], [2,2,0]\}$ forms a $(3,2)_3$ -PF (the sum of elements in each cycle is congruent to 1 mod 3).
- $D(S) = \{[0,1,2]\}$, is a $(3,1)_3$ -PF.

Justification I

Claim: If $\mathbf{a} = [a_0, a_1, \dots, a_{n-1}] \in S$, then $D^{-1}(D(\mathbf{a})) \subseteq S$.

Proof: First note that the elements of $D^{-1}(D(\mathbf{a}))$ must have period n , since $D(\mathbf{a})$ must have weight 0.

If $\mathbf{b} = [b_0, b_1, \dots, b_{n-1}] \in D^{-1}(D(\mathbf{a}))$, then for every i :
 $b_i = a_i + t$ for some t .

Hence:

$$b_0 + b_1 + \dots + b_{n-1} = a_0 + a_1 + \dots + a_{n-1} + nt \equiv 1 \pmod{c}$$

since $c|n$.

Justification II

- We next need to show that $|D^{-1}(D(\mathbf{a}))|=c$.
- Unfortunately, for $n=c=10$ this will not hold, since, for example,
 $D([0000355558])=D([5555800003])$.
- It will hold if we make the extra assumption that c is odd.

Justification III

- By the claim, the number of distinct cycles in $D(S)$ must be $|S|/n = c^{n-2}/n$.
- Hence the cycles in $D(S)$ contain a total of c^{n-2} $(n-2)$ -tuples.
- Since every $(n-1)$ -tuple occurs in a cycle in S , every $(n-2)$ -tuple must occur in a cycle in $D(S)$.
- Hence, every $(n-2)$ -tuple must occur in a unique cycle in $D(S)$.

Implication

- Unfortunately this construction does not give us a $(10,8)_{10}$ -PF, since here c is even!

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Finishing the job

- Over the last 20 years we have assembled a powerful set of construction techniques for de Bruijn graph perfect factors.
- This in turn allows us to construct examples of perfect maps for ‘most’ parameter sets.
- However, there is the fear that from now on we will just be knocking off sporadic cases.

A general approach?

- It seems possible that the construction shown is just a special case of a much more general ‘simple’ construction technique.
- I am hopeful that this can be used to cover many more previously undecided parameter sets.
- Fundamentally, perfect factors are very numerous.